



ALBUQUERQUE
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Introduction To Fortran 90/95

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Topics Covered Today

- ▶ History of Fortran standardization
 - ◆ **Why Fortran?**
- ▶ Overview of the language
 - ◆ **High level view**
 - ◆ **basics, procedures, modules, interfaces**



History

- ▶ Defined and implemented by Backus and his IBM team in mid 1950s
- ▶ Designed to demonstrate that a high level language for scientific computation could be efficient -- generate efficient object code
- ▶ Maintained its emphasis on efficiency as well as portability, particularly for numeric computation, throughout its development



History Contd

- ▶ Standardized in 1966 by ANSI
- ▶ First programming language to be standardized by an official body
- ▶ Revised 1978 -- Fortran 77 by ANSI
- ▶ Rev.1992 -- Fortran 90 by ANSI and ISO
- ▶ Revised 1997 -- Fortran 95 by ISO
- ▶ Technical development of Fortran 2K going on now
 - ◆ **completion expect in 2002 or so**



References And Textbooks

- ▶ The Fortran 95 Handbook, Adams, Brainerd, Martin, Smith, Wagener, MIT Press, 1997
- ▶ Fortran Top 90, Adams, Brainerd, Martin, Smith, Unicomp, 1994
- ▶ Programmer's Guide To Fortran 90, Brainerd, Goldberg, Adams, McGraw-Hill, 1990
- ▶ Fortran 90 For Scientists And Engineers, Nyhoff, Leestma, Prentice-Hall, 1997



What Is A Programming Language?

- ▶ A tool for instructing machines
- ▶ A means of communicating between programmers
- ▶ A vehicle for expressing high-level designs
- ▶ A notation for algorithms
- ▶ A way of expressing relationships between concepts -- a specification language or tool



A Programming Language?

- ▶ A tool for experimentation
- ▶ A means of controlling computerized devices
- ▶ A quote from Stroustrup, 1994
 - ◆ **NOT a collection of “neat” features**
 - ◆ **the workshop necessarily presents a collection of features but the features presented are driven from the applications**



What is Fortran's Niche?

- ▲ Efficiency of execution of the object code
 - ◆ **Because the standard's definition of the language has the following properties:**
 - extensive freedom to rearrange code provided the result is mathematically equivalent
 - the standard only specifies what the statements of the language mean -- not how they are to be implemented
 - no spurious constraints on how the language is to be implemented



Fortran's Niche Cont'd

- ▶ Portability of source code
 - ◆ **standardization**
 - ◆ **consistency with floating-point hardware standards**
- ▶ Major support for scientific (numer) comp
 - ◆ **extensive intrinsic library**
 - ◆ **extensive numeric precision control**
 - ◆ **full support of the complex data type**



Nu m e r i c P r e c i s i o n C o n t r o l

- ▶ Multiple floating point precisions
 - ◆ **single, double, and possible more**
 - ◆ **intrinsic math libraries for all precision provided**
 - ◆ **multiple precisions for the complex data type**
- ▶ Numeric inquiry and parameterization
 - ◆ **machine precision, tiny, huge, ...**
 - ◆ **consistent with IEEE binary floating-point standard 754**
- ▶ Manipulation of numeric parts of floating-point numbers
 - ◆ **fraction, exponent, scaling by the base, ...**



Fortran — An Imperative (Procedural) Language

- ▶ Other imperative languages
 - ◆ **C, Ada, Algols, Pascal, ...**
- ▶ Other kinds of languages
 - ◆ **functional languages (e.g. Lisp)**
 - ◆ **data flow languages (e.g. SISAL, Cantata)**
 - ◆ **object-oriented languages (e.g. Smalltalk, C++)**
 - ◆ **logic programming languages (e.g. Prolog)**
 - ◆ **specification languages (e.g. Ada, Z)**



Characteristic of Imperative Language

- ▶ Emphasis on executable statements
- ▶ Emphasis on data, types, values, operators
- ▶ Emphasis on procedures (functions and subroutines) as a means of extension
 - ◆ **to perform new operations or actions not intrinsic in the language**



Imperative Language Cont'd

- ▲ Emphasis on:
 - ◆ **efficiency**
 - ◆ **low level**
 - ◆ **systems and operations**
 - ◆ **imposes no discipline on interfaces between operations, procedures, parts**



Programming Vs Natural Languages

- ▶ Both are a means of communicating, human to machine, human to human
- ▶ Both have a structure:
 - ◆ a natural language has a grammar and a set of conventions or rules for expressing thoughts
 - ◆ a programming language has a **STRICT** grammar and set of rules
- ▶ Both have statements:
 - ◆ in a procedural programming language, they specify an action or specifies the environment in which the action takes place.



Comparison Continued

- ▶ Both have forms
 - ◆ **For a natural language**
 - books, chapters, paragraphs, sentences
 - ◆ **For a procedural language**
 - executable programs, program units, constructs, statements
- ▶ Both have rules of well-formed compositions
 - ◆ **For a natural language**
 - interpretation by context, ambiguities
 - ◆ **For a programming language**
 - strict adherence to rules of grammar and composition



An Example: DO Loops in Fortran 90

- ▶ It is a construct, that is, has some structure
 - ◆ a heading statement
 - ◆ a block of executable statements
 - ◆ a terminating statement
- ▶ Alternative forms of heading statement
 - ◆ do
 - ◆ do while(<exp>)
 - ◆ do <iteration_clause>
- ▶ Terminating statement (other forms possible)
 - ◆ enddo



DO Loops Continued

▶ An iterated DO Loop

```
sum = 0.0
```

```
do j = 1, n
```

```
    sum = sum + a(j)*b(j)
```

```
enddo
```

- ◆ initial value and final value of do-variable range can be any expression
- ◆ a stride can be specified as a third item
 - it may be negative but must not be zero
- ◆ any statement can be in the DO-loop body including other DO loops



DO Loops Continued

- ▶ A DO-while loop:

```
continue = .true.
```

```
do while ( continue )
```

```
    xnew = xold + step(xold)      ! step is a function
```

```
    if( abs(xnew-xold) < epsilon(xold) ) then
```

```
        continue = .false.
```

```
    else
```

```
        xold = xnew
```

```
    endif
```

```
enddo
```



DO Loops Continued

▶ The DO forever:

do

.... ! Before some computation

..... ! Check for convergence

if(converged) exit

..... ! Complete computation for next step

if(computation_ok) cycle

..... ! Find a better iterate and repeat

enddo



English Vs. Fortran

- | | |
|---------------------|---------------------|
| ▶ Book | Complete program |
| ▶ Chapters | Subprogram |
| ▶ Paragraphs | Parts of a program |
| ▶ Title of chapter | Subprogram header |
| ▶ Intro. paragraph | Specification part |
| ▶ Middle paragraphs | Execution part |
| ▶ Sections | Internal procedures |



English Vs. Fortran Cont'd

- ▶ Summary paragraph
- ▶ Sentences
 - ◆ **assertive**
 - ◆ **question**
- ▶ Words
 - ◆ **defined in dictionary**

End statement
Statements

- **spec. statements**
- **executable stmts**

Tokens

- **identifiers, operators**
- **keywords**



English Vs. Fortran Cont'd

▶ Delimiters

- ◆ **periods, commas, parenthesis, question marks, blanks**

▶ Parts of words

- ◆ **letters, apostrophe, digits**

▶ Side (margin) notes

Delimiters

- **periods, commas, parenthesis, ...**

Parts of tokens

- **letters, digits, underscores**

Comments



English Vs. Fortran Cont'd

▲ References

- ◆ **see ...**
- ◆
- ◆ **[ref]**

References

- **CALL statements**
- **function references**
- **Include statements**
- **USE statements**



Comparison Summary

- ◆ For a natural language, you learn rules for constructing sentences, paragraphs, sections, chapters, books, etc. and memorize meanings of words to construct the other constructs.
- ◆ In a programming language, you learn rules for constructing statements, subprograms, and programs, memorize semantics of individual statements, and construct and manipulate data objects with values to create a program.



An Example: hello.f90

- ▶ A simple Fortran program

```
program hello           ! Header statement

character(10) T        ! Specification part
intrinsic date_and_time

call date_and_time( time = T )  ! Exec. part
print *, 'Hello, world: the time is ', T(1:2), ":", &
        T(3:4)

end program hello      ! End statement
```



Points To Be Made About Helib.f90

- ▶ Program structure
 - ◆ header statement (optional for main program)
 - ◆ specification part (zero or more statements)
 - ◆ executable part (zero or more statements)
 - ◆ internal procedure part (zero or more statements)
 - ◆ end statement
- ▶ ALL program units have this structure
 - ◆ Program units include:
 - main program, subroutine subprogram, function subprogram, module, ..., internal procedure, module procedure



Statements From `help.f90`

- ▶ Header for a main program
 - ◆ **program <name_of_main_program>**
- ▶ Specification part
 - ◆ **type declaration CHARACTER(<length>) <identifiers>**
 - defines identifiers to be character strings of a specified length
 - ◆ **intrinsic statement INTRINSIC <names>**
 - specifies the names are intrinsic procedures
 - ◆ **many other kinds of specification statements**
 - **SAVE, POINTER, ALLOCATABLE, TYPE, DATA, COMMON, INTERFACE** blocks, ...



Statements From Helib.f90 Continued

- ▶ Executable part
 - ◆ **CALL statement -- CALL <name>(<arguments>)**
 - go to <name> (somewhere else) to execute code and return results, usually as arguments
 - ◆ **PRINT statement -- PRINT <format>, <list_of_output_items>**
 - print the values of variables listed in the I/O list with specified format
 - * for <format> means the processor selects a suitable form for printing the listed output items
 - output items can be variables, constants or expressions



Statements From Fortran 90 Continued

- ▶ Executable part continued
 - ◆ other kinds of statements than can be in the executable part
 - assignment statements
 - control constructs(DO loops, IF construct, CASE construct)
 - allocate/deallocate statements
 - special array statements (FORALL, WHERE)
 - I/O statements (READ, WRITE, OPEN, CLOSE, ...)
 - pointer-assignment statement
 - GOTO, CONTINUE, EXIT, CYCLE, ...

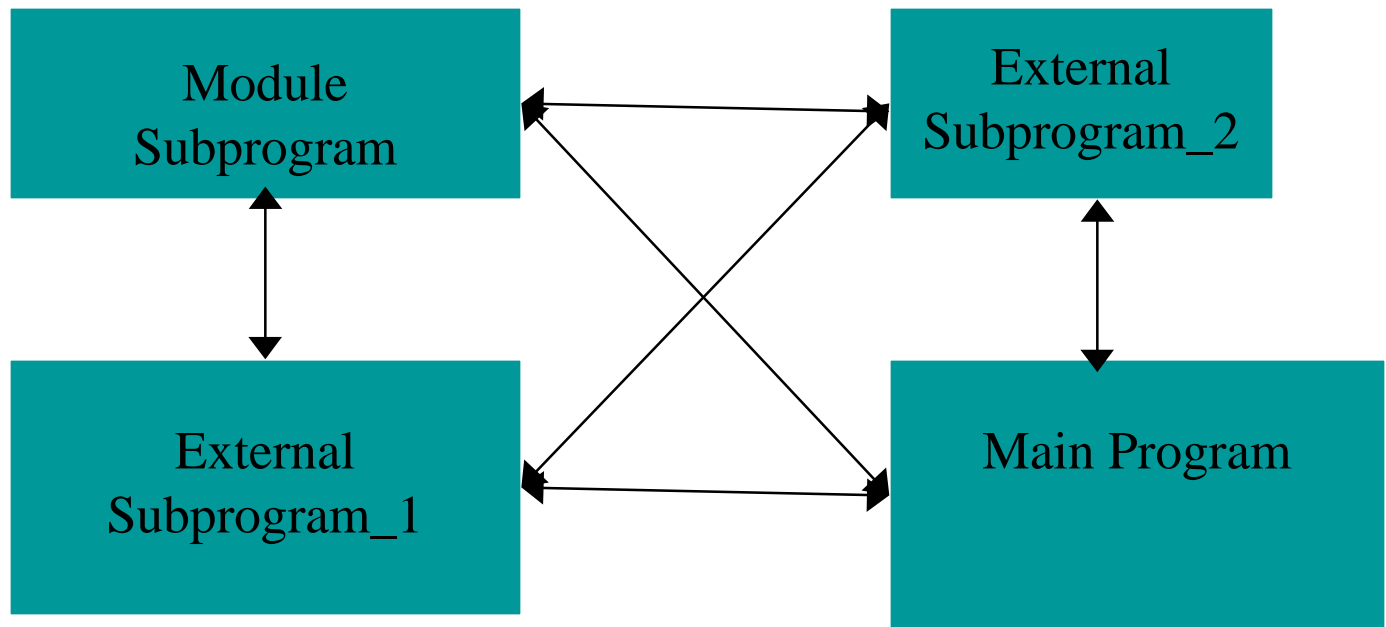


Statements From Hello.f90 Continued

- ▶ Internal procedure part
 - ◆ no internal procedures in Hello.f90
 - ◆ could be -- begins with a **CONTAINS** statement
 - followed by a subroutine or function program unit
- ▶ END statement
 - ◆ the word **END** followed optionally by:
 - 1) the word **PROGRAM** for the main program
 - SUBROUTINE for a subroutine
 - FUNCTION for a function
 - 2) the name of the program -- hello in this case



Executable Program — The Whole View



Arrows \leftrightarrow indicate association of names



Discussion of The Whole View

- ▶ Program execution starts with the main program
 - ◆ at some point, it calls an external subprogram
 - ◆ at some point, it calls a procedure in a module
 - ◆ it makes use of a variable in a module
- ▶ The connection between the main program and procedures or variables is via an association (the lines and arrows)
 - ◆ argument association
 - ◆ storage association
 - ◆ USE association



Why Is It Association?

- ▶ Most names are local to a program unit in Fortran
 - ◆ **For example, in a declaration in the main program**

```
character(10) T
```

- **T is a local variable in the main program and is NOT the same as any T in general in any other program unit**
- ▶ Global names
 - ◆ **external subprogram names, main program names, module names, common block names, and a very few other items**



Why Is It Association? Continued

- ▲ How does a local variable TIME in a function, say, take the value of a local T in the main program?
 - ◆ **Argument association**
 - while the CALL is being execution, T is the same as TIME
`CALL date_and_time(TIME = T)`
 - when the CALL completes, the association terminates
 - ◆ **Similarly with storage association, USE association, HOST association, etc. which we will learn about**



Separate Compilation — The Whole View

- ▲ Fortran permits (and assumes) the concept of separate compilation
 - ◆ external program units can be in separate file and can be compiled separately from parts of the program in other files
 - ◆ these means that the compiler has to **infer** what is expected by the called procedure from the call site
 - that is, the “interface” to an external program is implicit or guessed from the call site: For example, consider:

```
call subr(j)
```

 - If j is an integer, the compiler **assumes** the dummy argument is an integer.



An Example With Functions: simple_ext.f90

```
program test_function_simple
  implicit none
  real, parameter :: SMALL = 1.0e-15
  real, parameter :: MODEST = 2.0
  real, parameter :: LARGE = 2.0e15
  integer i
  real result
  integer, parameter :: N = 3
  real, dimension(N) :: x
  data x/SMALL, MODEST, LARGE/
  real, external :: simple ! Necessary because of impl. none
```



simple.f90 Continued

```
print *, '      x      simple(x)'  
  
do i = 1, N  
    result = simple(x(i))  
    print *, x(i), result  
enddo  
  
end program test_function
```



simple.f90 Continued

```
function simple( x )  
  
    implicit none  
    real x  
    real simple  
    intrinsic sqrt  
  
    simple = sqrt( 1.0 + X**2 )  
  
end function simple
```



Points About simple.f90

- ▶ The previous three slides are in one file
 - ◆ **Two program units:**
 - main program plus a function subprogram
 - both are external units
 - in particular, an external function
 - ◆ **The two units may be in separate files**
 - separate and independent compilation units
 - ◆ **Communication between units:**
 - **argument association only**
 - $x(i)$ in main is associated with dummy x in function, only while the call is being executed
 - two way association -- both in and out



Points About simple.f90 Continued

- ▲ Details about the program
 - ◆ **function reference in an expression**
 - the right hand side of an assignment statement
 - this expression is simple here but can be arbitrary expression
 - ◆ **inside the function subprogram, the name simple is both**
 - the place to define the result as a local variable
 - the name of the global function known in other program units
 - a result clause can be used to distinguish the two uses

```
function simple( x ) result( res )  
res = sqrt( 1.0 + x**2 )
```




Points About simple.f90 Continued

- ▶ Declarations:
 - ◆ **PARAMETER attribute**
 - defines a named constant -- gives a name to a constant value
 - ◆ **Attributes can be grouped together in a type statement**
`<type> [, <list_of_attributes> ::] <list_of_identifiers>`
 - ◆ **Examples:**

```
real, parameter :: modest = 2.0  
real, external :: simple
```
 - ◆ **The declaration of the name simple is required because implicit none means all identifiers have to be specified**



Interfaces

- ▶ The declaration of the identifier simple is providing a partial interface to a procedure
 - ◆ **An interface consists of:**
 - the name and number of arguments, and return type
 - the type, name, and other attributes of its dummy arguments
 - other attributes are:
 - dimension, pointer, optionality, intent, ...
- ▶ This interface is partial because:
 - ◆ **only types of returned result and argument are known to the calling program unit**



Interfaces Continued

- ▲ The complete interface to an external program can be provided by an interface specification block
 - ◆ appears in the calling program in the specification part
 - ◆ its form is the header, argument specification, and end statement of a subprogram unit

```
interface
  function simple( x )
    real x, simple
  end function
end interface
```



Interfaces Continued

- ▶ The interface specification makes the interface explicit for external subprogram
 - ◆ **otherwise, the interface is implicit**
 - has to be guessed or assumed by the compiler
- ▶ There are other kinds of procedures
 - ◆ **internal and module procedures**
 - internal procedures are those defined **INSIDE** a main or procedure program unit
 - module procedures are those defined **INSIDE** a module
- ▶ These procedures also have explicit interfaces because they are visible



Simple.f90 Using An InternalFunction

- ▶ See file simple_int.f90
 - ◆ **The code looks like:**

Main Program

Code for main program

CONTAINS

Function simple
(all code)

End Main Program



Modules

- ▶ Modules are program units that contain:
 - variables (data)
 - constants (fixed data)
 - other kinds of data objects
 - interfaces
 - procedures
 - ◆ that can be referenced and used by other program units
- ▶ These items are referenced as follows:
 - ◆ a **USE** statement in any program that wants to use them
 - ◆ they are known by their name
 - ◆ a module provides “selective global entities”



Modules Continued

- ▶ “Selective global entities:
 - ◆ known **only** to programs that reference the module by a **USE** statement
 - ◆ never conflict with local names in program unit that do not reference the module with a **USE** statement
 - ◆ In other words, there can be many definitions of the module function simple in the same executable file
- ▶ Permits data abstraction (and data hiding)
 - ◆ can control access to data and access to procedures (using **PRIVATE** and **PUBLIC** specifications)



Simple.f90 Using A Module Function

- See file simple_mod.f90

```
Module my_functions  
Module data goes here (if any)  
CONTAINS  
    Function simple  
        (all code)  
End module my_functions
```

```
Main Program  
    USE my_functions  
    Code for main program  
End main program
```




A Numerically Careful Implementation Of Simple.f90 — Computes $\sqrt{1+x^2}$

- ▶ As written, simple.f90 is not very robust
 - ◆ for large x , it may overflow spuriously and either
 - stop execution
 - continue with an overflow message
 - ◆ for small x , it may underflow spuriously and either
 - continue with no message
 - continue with a message
 - stop with a diagnostic
 - ◆ Why is there really no problem?
 - If x is large in magnitude relative to 1, the answer is x
 - If x is small in magnitude relative to 1, the answer is 1



Careful Implementation Continued

- ▶ How do we know what is small or large relative to 1
 - ◆ **depends on the precision of the arithmetic used**
 - **suppose 7 digits of precision we used**
 - then, if $|x| > 10^7$, square root of $1 + x^2$ is x within 1 rounding error
 - if $|x| < 10^{-7}$, square root of $1 + x^2$ is 1 within 1 rounding error
 - **TO BE COMPLETED**